

This research contains as a main result the proof that every Chordal B1-EPG simultaneously in the VPT and EPT graph classes. In graph **1S** addition, we describe a set of graphs that defines Helly- B_1 -EPG families. In particular, this work presents some features of non-trivial families of graphs properly contained in Helly-B₁ EPG, namely Bipartite, Block, Cactus and Line of Bipartite graphs.

In this work we will mainly explore the EPG graphs, in particular B_1 -EPG graphs. However, other classes of intersection graphs will be studied such as EPT and VPT graph classes.

- has at most k bends;
- EPG representation;
- In a B₁-EPG representation, a clique K can be **edge-clique** or **claw-clique** [3].



Figure 1: Representation of a clique as edge-clique and as claw-clique.

- \bullet sub-collection has at least one common element;
- a Helly representation;
- Helly-B₁-EPG graphs were studied in [2];
- imization [4];
- \bullet paths on trees, respectively;
- VPT and EPT graphs are incomparable families of graphs.

Paths On Hosts: B₁-EPG, EPT and VPT Graphs

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Introduction

Objective

Definitions and Technical Results

A graph is a B_k -EPG graph if it admits an EPG representation in which each path [•]

When k = 1 we say that this is a single bend EPG representation or simply a B₁-



A collection of sets satisfies the **Helly property** when every pairwise intersecting

When this property is satisfied by the set of paths used in a representation, we get

EPG, EPT and VPT representations arise in circuit layout problems and layout opt

VPT and EPT graphs are the vertex-intersection and edge-intersection graphs of

Subclasses of Helly-B1-EPG Graphs

Theorem 1: Let G be a B_1 -EPG graph. If G is $\{S_3, S_3, S_3, S_3, C_4\}$ -free then G is a Helly- B_1 -EPG graph.



(a) Claw with paths.



(b) Subgraph induced by paths.





Figure 3: Graphs on statement of Theorem 1: S₃, S_{3'}, S_{3''}, C₄.

Bull-free graphs are $\{S_3, S_{3'}, S_{3''}\}$ -free, so these results implies in results of [1].

Theorem 2: If the graph G is B_1 -EPG and diamond-free then G is Helly- B_1 -EPG. Corollary: Bipartite, Block, Cactus and Line of Bipartite graphs are Helly-B1-EPG.

Relationship among Chordal B1-EPG, VPT and EPT graphs

Theorem 3: Chordal B_1 *-EPG* \subseteq *VPT*.







Figure 4: Graph S₄ and one of its possible VPT and EPT representations.

References

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Theorem 4: Chordal B_1 *-EPG* \subseteq *EPT*.

